The Forcing Connected Weak Edge Detour Number of a Graph

J.M. Prabakar, I. Keerthi Asir and S. Athisayanathan

Department of Mathematics St. Xavier's College (Autonomous), Palayamkottai - 627 002 Tamilnadu, India.

jmpsxc@gmail.com, as irsxc@gmail.com, at hisxc@gmail.com

Abstract

In this paper, we define the forcing connected weak edge detour set (number $(fcdn_w(G))$, basis). Also, the forcing connected weak edge detour number of certain standard of graphs are determined. It is shown that for each pair a, b of integers with $0 \le a \le b$ and $b \ge 3$, there is a connected graph G with $fcdn_w(G) = a$ and $cdn_w(G) = b$, where $cdn_w(G)$ is connected weak edge detour number of G.

Key words: Connected weak edge detour number, forcing weak edge detour number, forcing connected weak edge detour number.

2010 Mathematics Subject Classification : 05C40, 05C69

1 Introduction

By a graph G = (V, E), we mean a finite undirected connected simple graph. The order and size of G are denoted by n and m respectively. For basic definitions and terminologies, we refer to [1, 4].

For vertices u and v in a connected graph G, the distance d(u, v) is the length of a shortest u - v path in G. A u - v path of length d(u, v) is called a u - v geodesic. For a vertex v of G, the eccentricity e(v) is the distance between v and a vertex farthest from v. The minimum eccentricity among the vertices of G is the radius, rad (G) and the maximum eccentricity is its diameter, diam (G) of G.

For vertices u and v in a connected graph G, the detour distance D(u, v) is the length of a longest u - v path in G. A u - v path of length D(u, v) is called a u - v detour. The detour eccentricity $e_D(v)$ of a vertex v in G is the maximum detour distance from v

^{*} Corresponding Author: J.M. Prabakar

 $[\]Psi$ Received on July 12, 2017 / Revised on October 01, 2017 / Accepted on October 13, 2017

to a vertex of G. The detour radius, $rad_D(G)$ of G is the minimum detour eccentricity among the vertices of G, while the detour diameter, $diam_D(G)$ of G is the maximum detour eccentricity among the vertices of G. A vertex x is said to lie on a u-v detour Pif x is a vertex of P including the vertices u and v. A set $S \subseteq V$ is called a detour set if every vertex v in G lies on a detour joining a pair of vertices of S. The detour number dn(G) of G is the minimum order of a detour set and any detour set of order dn(G) is called a detour basis of G. These concepts were studied by Chartrand et al.[3].

A subset $T \subseteq S$ is called a *forcing subset* for S if S is the unique minimum detour set containing T. A forcing subset for S of minimum cardinality is a *minimum forcing subset* of S. The *forcing detour number* of S, denoted by fdn(S), is the cardinality of a minimum forcing subset for S. The forcing detour number of S, denoted by fdn(S), is the cardinality of forcing subset for S. The forcing detour number of G is fdn(G) = $min\{fdn(S)\}$, where the minimum is taken over all minimum connected weak edge detour sets S in G. The forcing detour number of a graph was introduced and studied by Chartrand et.al [3].

Let G be a connected graph and S a connected detour basis of G. A subset $T \subseteq S$ is called a *forcing subset* for S if S is the unique minimum detour set containing T. A forcing subset for S of minimum cardinality is a *minimum forcing subset* of S. The *forcing connected detour number* of S, denoted by fcdn(S), is the cardinality of a minimum forcing subset for S. The forcing detour number of S, denoted by fcdn(S), is the cardinality of a minimum forcing subset for S. The forcing detour number of S, denoted by fcdn(S), is the cardinality of forcing subset for S. The forcing detour number of G is $fcdn(G) = min\{fcdn(S)\}$, where the minimum is taken over all minimum connected detour sets S in G. The forcing connected detour number was introduced and studied by Santhakumaran and Athisayanathan [5].

Let G be a connected graph and S a connected weak detour basis of G. A subset $T \subseteq S$ is called a *forcing subset* for S if S is the unique minimum detour set containing T. A forcing subset for S of minimum cardinality is a *minimum forcing subset* of S. The *forcing weak edge detour number* of S, denoted by fcdn(S), is the cardinality of a minimum forcing subset for S. The forcing weak detour number of S, denoted by $fdn_w(S)$, is the cardinality of forcing subset for S. The forcing weak detour number of G is $fdn_w(G) = min\{fdn_w(S)\}$, where the minimum is taken over all minimum weak edge detour sets S in G. In 2010, the forcing weak edge detour number of a graph was introduced and studied by Santhakumaran and Athisayanathan [6].

A set $S \subseteq V$ is called a *connected weak edge detour set* of a graph G, if S is a weak edge

detour set and the subgraph $\langle S \rangle$ induced by S is connected. The minimum cardinality of a connected weak edge detour set of G is connected weak edge detour number, denoted by $cdn_w(G)$ of G and any connected weak edge detour set of order $cdn_w(G)$ is called a connected weak edge detour basis of G. These concepts were studied by Prabakar and Athisayanathan [7].

The following theorems are used in sequel.

Theorem 1.1. [7] Let G be the complete graph K_n $(n \ge 2)$. Then a set $S \subseteq V$ is a connected weak edge detour basis of G if and only if S consists of any two vertices of G.

Theorem 1.2. [7] Let G be the complete bipartite graph $K_{m,n}$ $(2 \le m \le n)$. Then a set $S \subseteq V$ is a connected weak edge detour basis of G if and only if S consists of any two adjacent vertices of G.

Theorem 1.3. [7] Let G be a cycle of order $n \ge 3$. Then a set $S \subseteq V$ is a connected weak edge detour basis of G if and only if S consists of any two adjacent vertices of G.

Theorem 1.4. [7] For any graph G of order $n \ge 2, 2 \le cdn_w(G) \le n$.

Theorem 1.5. [7] Let $G = (K_{n_1} \cup K_{n_2} \cup \ldots \cup K_{n_r} \cup kK_1) + v$ be a block graph of order $n \ge 4$ such that $r \ge 1$, each $n_i \ge 2$ and $n_1 + n_2 + \ldots + n_r + k = n - 1$. Then $cdn_w(G) = r + k + 1$.

Theorem 1.6. [7] For any connected graph G with k end-vertices and l cut-vertices, $max\{2, k+l\} \leq cdn_w(G) \leq n.$

Throughout this paper G denotes a connected graph with at least two vertices.

2 Forcing subsets in Connected Weak Edge Detour Sets

Even though every connected graph contains a connected weak edge detour basis, some connected graph may contain several connected weak edge detour bases. For each connected weak edge detour basis S in a connected graph G, there is always some subset T of S that uniquely determines S as the connected weak edge detour basis containing T. Such "forcing subsets" will be considered in this section.

Definition 2.1. Let G be a connected graph and S a connected weak edge detour basis of G. A subset $T \subseteq S$ is called a *forcing subset* for S if S is the unique connected weak edge detour basis containing T. A forcing subset for S of minimum cardinality is a *minimum*

forcing subset of S. The forcing connected weak edge detour number of S, denoted by $fcdn_w(S)$, is the cardinality of a minimum forcing subset for S. The forcing connected weak edge detour number of G, denoted by $fcdn_w(G)$, is $fcdn_w(G) = min\{fcdn_w(S)\}$, where the minimum is taken over all minimum connected weak edge detour bases S in G.

Example 2.2. For the graph G given in Figure 2.1, $S = \{u, v, w, x, y\}$ is the unique connected weak edge detour basis of G so that $fcdn_w(G) = 0$. For the graph G given in Figure 2.2, $S_1 = \{u, s, w, t, v\}$, $S_2 = \{u, s, x, t, v\}$ and $S_3 = \{u, s, y, t, v\}$ are the three connected weak edge detour bases of G so that $fcdn_w(G) = 1$.



Figure 2.1



Example 2.3. For the graph G given in Figure 2.3, fcdn(G) = 2 and $fcdn_w(G) = 1$. Hence the forcing connected detour number and forcing connected weak edge detour number are different.



Theorem 2.4. For every connected graph $G, 0 \leq fcdn_w(G) \leq cdn_w(G)$.

Proof: It is clear from the definition of $fcdn_w(G)$ that $fcdn_w(G) \ge 0$. Let S be any connected weak edge detour basis of G. Since $fcdn_w(S) \le cdn_w(G)$ and since $fcdn_w(G) = min\{fcdn_w(S); S \text{ is a connected weak edge detour basis of } G\}$, it follows that $fcdn_w(G) \le cdn_w(G)$. Thus $0 \le fcd_n(G) \le cdn_w(G)$.

Remark 2.5. The lower bound in theorem 2.4 are sharp. For the graph G given in Figure 2.1, $fcdn_w(G) = 0$. For the cycle C_3 , $fcdn_w(C_3) = cdn_w(C_3) = 2$. Also, all the inequalities in theorem 2.4 can be strict. For the graph G in Figure 2.2, $fcdn_w(G) = 1$ and $cdn_w(G) = 5$ so that $0 < fcdn_w(G) < cdn_w(G)$.

In the following theorem we characterize graphs G for which the bounds in theorem 2.4 are attained and also graphs for which $fcdn_w(G) = 1$.

Theorem 2.6. Let G be a connected graph. Then

- (a) $fcdn_w(G) = 0$ if and only if G has a unique connected weak edge detour basis.
- (b) $fcdn_w(G) = 1$ if and only if G has at least two connected weak edge detour bases, one of which is a unique connected weak edge detour basis containing one of its elements, and
- (c) $fcdn_w(G) = cdn_w(G)$ if and only if no connected weak edge detour basis of G is the unique connected weak edge detour basis containing any of its proper subsets.
- **Proof:** (a) Let $fcdn_w(G) = 0$. Then, by definition, $fcdn_w(S) = 0$ for some connected weak edge detour basis S of G so that the empty set ϕ is the minimum forcing subset for S. Since the empty set ϕ is a subset of every set, it follows that S is the unique connected weak edge detour basis of G. The converse is clear.
 - (b) Let $fcdn_w(G) = 1$. Then by (a), G has at least two connected weak edge detour bases. Also, since $fcdn_w(G) = 1$, there is a singleton subset T of a connected weak edge detour basis S of G such that T is not a subset of any other connected weak edge detour basis of G. Then S is the unique connected weak edge detour basis containing one of its elements. The converse is clear.

(c) Let $fcdn_w(G) = cdn_w(G)$. Then $fcdn_w(S) = cdn_w(G)$, for every connected weak edge detour basis S in G. Also, by theorem 1.4, $cdn_w(G) \ge 2$ and hence $fcdn_w(G) \ge 2$. Then by (a), G has at least two connected weak edge detour bases and so the empty set ϕ is not a forcing subset of any weak edge detour basis of G. Since $fcdn_w(S) = cdn_w(G)$, no proper subset of S is a forcing subset of S. Thus no connected weak edge detour basis of G is the unique connected weak edge detour basis containing any of its proper subsets. Conversely, the data implies that Gcontains more than one connected weak edge detour basis and no subset of any connected weak edge detour basis S other than S is a forcing subset for S. Hence it follows that $fcdn_w(G) = cdn_w(G)$.

Theorem 2.7. Let G be a connected graph and let \mathcal{F} be the set of relative complements of the minimum forcing subsets in their respective connected weak edge detour basis in G. Then $\bigcap_{F \in \mathcal{F}} F$ is the set of connected weak edge detour edges of G.

Proof: Let W be the set of all connected weak edge detour vertices of G. We claim that $W = \bigcap_{F \in \mathcal{F}} F$. Let a vertex $v \in W$. Then v is a connected weak edge detour vertex of G so that v belongs to every connected weak edge detour basis S of G. Let $T \subseteq S$ be any minimum forcing subset for any connected weak edge detour basis S of G. We claim that $v \notin T$. If $v \in T$, then $T' = T - \{v\}$ is a proper subset of T such that S is the unique connected weak edge detour basis containing T' so that T' is a forcing subset for S with $|T'| \leq |T|$, which is a contradiction to T is a minimum forcing subset for S. Thus $v \notin T$ and so $v \in F$, where F is the relative complement of T in S. Hence $v \in \bigcap_{F \in \mathcal{F}} F$ so that $W \subseteq \bigcap_{F \in \mathcal{F}} F$. Conversely, let $v \in \bigcap_{F \in \mathcal{F}} F$. Then v belongs to the relative complement of T in S for every T and every S such that $T \subseteq S$, where T is minimum forcing subset for S. Since F is the relative complement of T in S, $F \subseteq S$ and so $v \in S$, for every S so that v is a connected weak edge detour vertex of G. Thus $v \notin W$ and so $\bigcap_{F \in \mathcal{F}} F \subseteq W$. Hence $W = \bigcap_{F \in \mathcal{F}} F$.

Corollary 2.8. Let G be a connected graph and S a connected weak edge detour basis of G. Then no connected weak edge detour vertex of G belong to minimum forcing set of G.

Proof: The proof is contained in the first part of theorem 2.7.

Theorem 2.9. Let G be a connected graph and W be the set of all connected weak edge detour vertices of G. Then $fcdn_w(G) \leq cdn_w(G) - |W|$.

Proof: Let S be any connected weak edge detour basis of G. Then $cdn_w(G) = |S|$, $W \subseteq S$ and S is the unique connected weak edge detour basis containing S - W. Thus $fcdn_w(G) \leq |S - W| = |S| - |W| = cdn_w(G) - |W|$.

Corollary 2.10. if G is a connected graph with k end-vertices and l cut-vertices, then $fcdn_w(G) \leq cdn_w(G) - k - l$.

Proof: This follows from theorem 1.6 and theorem 2.9.

Remark 2.11. The bound in theorem 2.9 is sharp. For the graph G given in the Figure 2.4 the set $S_1 = \{v_1, v_2, v_3\}$, $S_2 = \{v_1, v_2, v_5\}$ are the connected detour bases of G so that $cdn_w(G) = 3$, |W| = 2 and $fcdn_w(G) = 1$ as in the Remark 2.5. Also, the inequality in theorem 2.9 can be strict. For the graph G given in Figure 2.5 the sets $S_1 = \{v_1, v_2, v_3\}$, $S_2 = \{v_1, v_2, v_4\}$ and $S_3 = \{v_1, v_3, v_5\}$ are three connected detour bases of G so that |W| = 1, $cdn_w(G) = 3$ and $fcdn_w(G) = 1$. Thus $fcdn_w(G) < cdn_w(G) - |W|$.

Moreover, the bound in Corollary 2.10 is also sharp. For the graph G given in Figure 2.4, $cdn_w(G) = 3$, k = 1, l = 1 and $fcdn_w(G) = 1$. Also, the inequality in Corollary 2.10 can be strict. For the graph G given in Figure 2.2, $cdn_w(G) = 5$, k = 1, l = 1 and $fcdn_w(G) = 1$. Thus $fcdn_w(G) < cdn_w(G) - k - l$.





- **Theorem 2.12.** (a) If G is the complete graph K_n $(n \ge 3)$ or the complete bipartite graph $K_{m,n}$ $(2 \le m \le n)$, then $cdn_w(G) = fcdn_w(G) = 2$.
 - (b) If G is the cycle C_n $(n \ge 3)$, then $cdn_w(G) = fcdn_w(G) = 2$.
 - (c) If G is a tree of order $n \ge 2$, then $cdn_w(G) = fcdn_w(G) = 0$.
- **Proof:** (a) By theorem 1.1 and theorem 1.2, a set S of vertices is a connected weak edge detour basis of G if and only if S consists of any two vertices of G. For each vertex v in G there are two or more vertices adjacent with v. Thus the vertex v belongs to more than one connected weak edge detour basis of G. Hence it follows that no set consisting of a single vertex is a forcing subset for any connected weak edge detour basis of G. Thus the result follows.
 - (b) By theorems 1.3 (according as G is odd or even), a set S of two adjacent vertices of G is a connected weak edge detour basis. For each vertex v in G there are two vertices adjacent with v. Thus the vertex v belongs to more than one connected weak edge detour basis of G. Hence it follows that no set consisting of a single vertex is a forcing subset for any connected weak edge detour basis of G. Thus the result follows.
 - (c) By theorem 1.4, $cdn_w(G) = n$. Since the set of all vertices of a tree is the unique connected weak edge detour basis so that $fcdn_w(G) = 0$ by theorem 2.6(a).

The following theorem gives a realization result.

Theorem 2.13. For each pair a, b of integers with $0 \le a \le b$ and $b \ge 3$, there is a connected graph G with $fcdn_w(G) = a$ and $cdn_w(G) = b$.

Proof: Case 1 : a = 0. For each $b \ge 3$, let G be a tree with b verifices. Then $fcdn_w(G) = 0$ and $cdn_w(G) = b$ by theorem 2.12(c).

Case 2 : $a \ge 1$. For each *i* with $(1 \le i \le a)$, let F_i be a copy of the complete graph K_2 , where $V(F_i) = \{u_i, v_i\}$ and let $H = K_{1,b-a-1}$ be the star at *v* whose vertex set is $W = \{z_1, z_2, \ldots, z_{b-a-1}, v\}$. Then the graph *G* is obtained by joining the central vertex *v* of *H* to the vertices of F_1, F_2, \ldots, F_a . The graph *G* is connected and is shown in Figure 2.6. Then by theorem 1.5, $cdn_w(G)=b$. Now, we show that $fcdn_w(G) = a$. It is clear that *W* is the set of all connected weak edge detour vertices of *G*. Hence it follows from Theorem 2.9 that $fcdn_w(G) \le cdn_w(G) - |W| = b - (b-a) = a$. Now, since $cdn_w(G) = b$, it follows from Theorem 1.11 that any connected weak edge detour basis of *G* is of the form $S = W \cup \{x_1, x_2, \ldots, x_a\}$, where $x_i \in \{u_i, v_i\}$ $(1 \le i \le a)$. Let *T* be a subset of *S* with |T| < a. Then there is a vertex x_j $(1 \le j \le a)$ such that $x_j \notin T$. Let y_j be a vertex of F_j distinct from x_j . Then $S' = (S - \{x_j\}) \cup \{y_j\}$ is also a connected weak edge detour basis containing *T* so *T* is not a forcing set of *S*. Since this is true for all connected weak edge detour basis detour bases of *G*, it follows that $fcdn_w(G) \ge a$ and so $fcdn_w(G) = a$.



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